

# an Introduction to Quantum Information Theory

**Claude Crépeau**

School of Computer Science  
McGill University



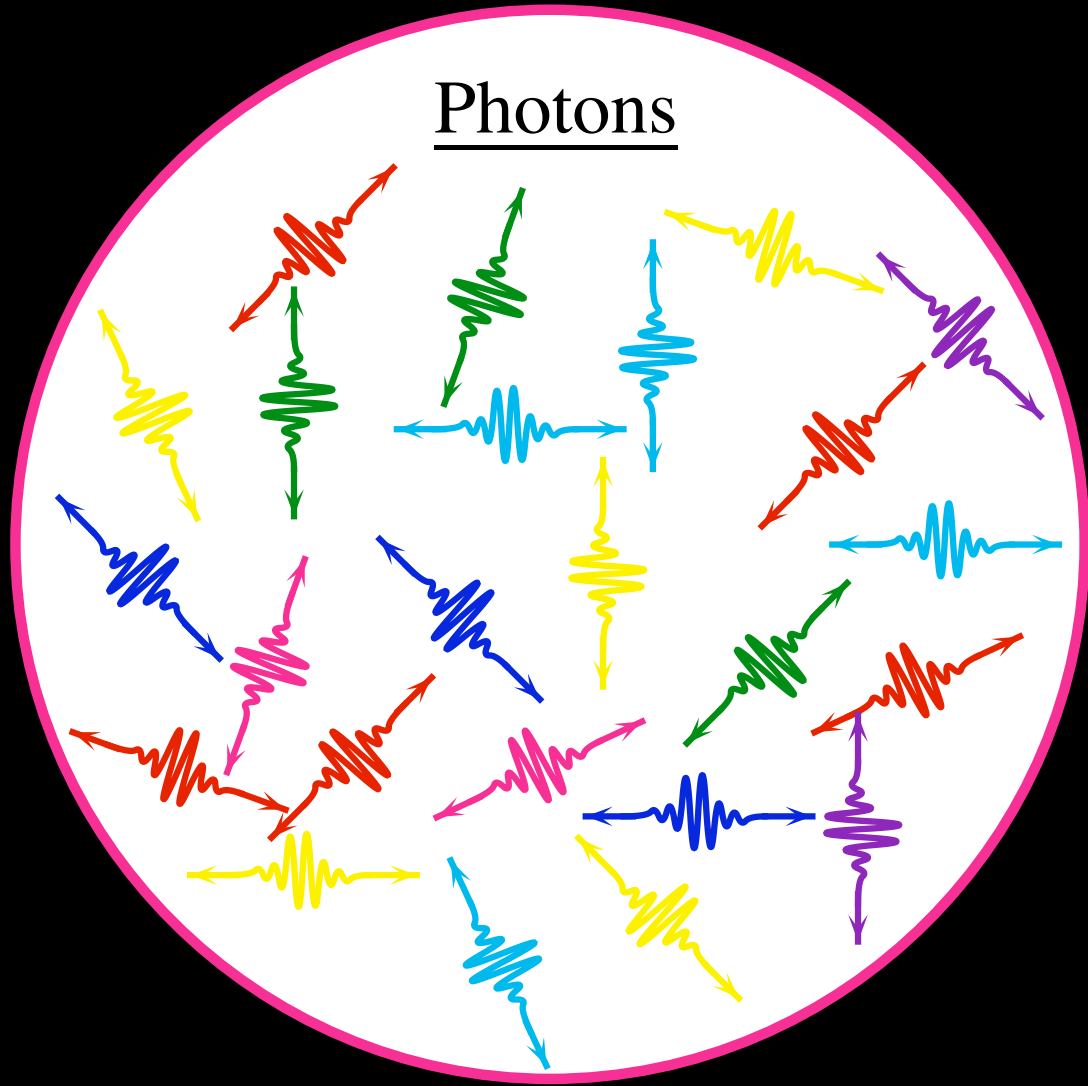
**(1)**

# **Quantum Information**

# Photons



# Photons



# Photons

# Bits & QuBits

0: 

1: 

$$\theta = \cos\theta \left\langle \leftarrow \right\rangle + \sin\theta \left\langle \updownarrow \right\rangle$$


$$|\Psi\rangle = C_0 \left\langle \leftarrow \right\rangle + C_1 \left\langle \updownarrow \right\rangle$$

$C_i, C_{ij} \in \mathbb{C}$

00: 

01: 

10: 

11: 

$$|\Psi\rangle = C_{00} \left\langle \leftarrow \right\rangle \left\langle \leftarrow \right\rangle + C_{01} \left\langle \leftarrow \right\rangle \left\langle \updownarrow \right\rangle + C_{10} \left\langle \updownarrow \right\rangle \left\langle \leftarrow \right\rangle + C_{11} \left\langle \updownarrow \right\rangle \left\langle \updownarrow \right\rangle$$

## standard notations

Basis vectors:  $|0\rangle$  and  $|1\rangle$

Arbitrary states:  $|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$   
such that  $|\alpha|^2 + |\beta|^2 = 1$

Arbitrary multi-states:

$$\alpha, \beta, \delta, \gamma \in \mathbb{C}$$

$|\psi\rangle = \alpha|00\rangle + \beta|01\rangle + \delta|10\rangle + \gamma|11\rangle$   
such that  $|\alpha|^2 + |\beta|^2 + |\delta|^2 + |\gamma|^2 = 1$

## standard notations

Conjugate: if  $|\psi\rangle = \alpha|0\rangle + \beta|1\rangle = \begin{pmatrix} \alpha \\ \beta \end{pmatrix}$

then  $\langle\psi| = \alpha^*\langle 0| + \beta^*\langle 1| = (\alpha^* \ \beta^*)$

Scalar product:  $\langle\psi|\phi\rangle = \langle\psi| \bullet |\phi\rangle$

$\langle 0|0\rangle = \langle 1|1\rangle = 1$  and  $\langle 0|1\rangle = \langle 1|0\rangle = 0$

Tensor product:  $|\phi\rangle \times |\psi\rangle = |\phi\rangle \bullet \langle\psi|$

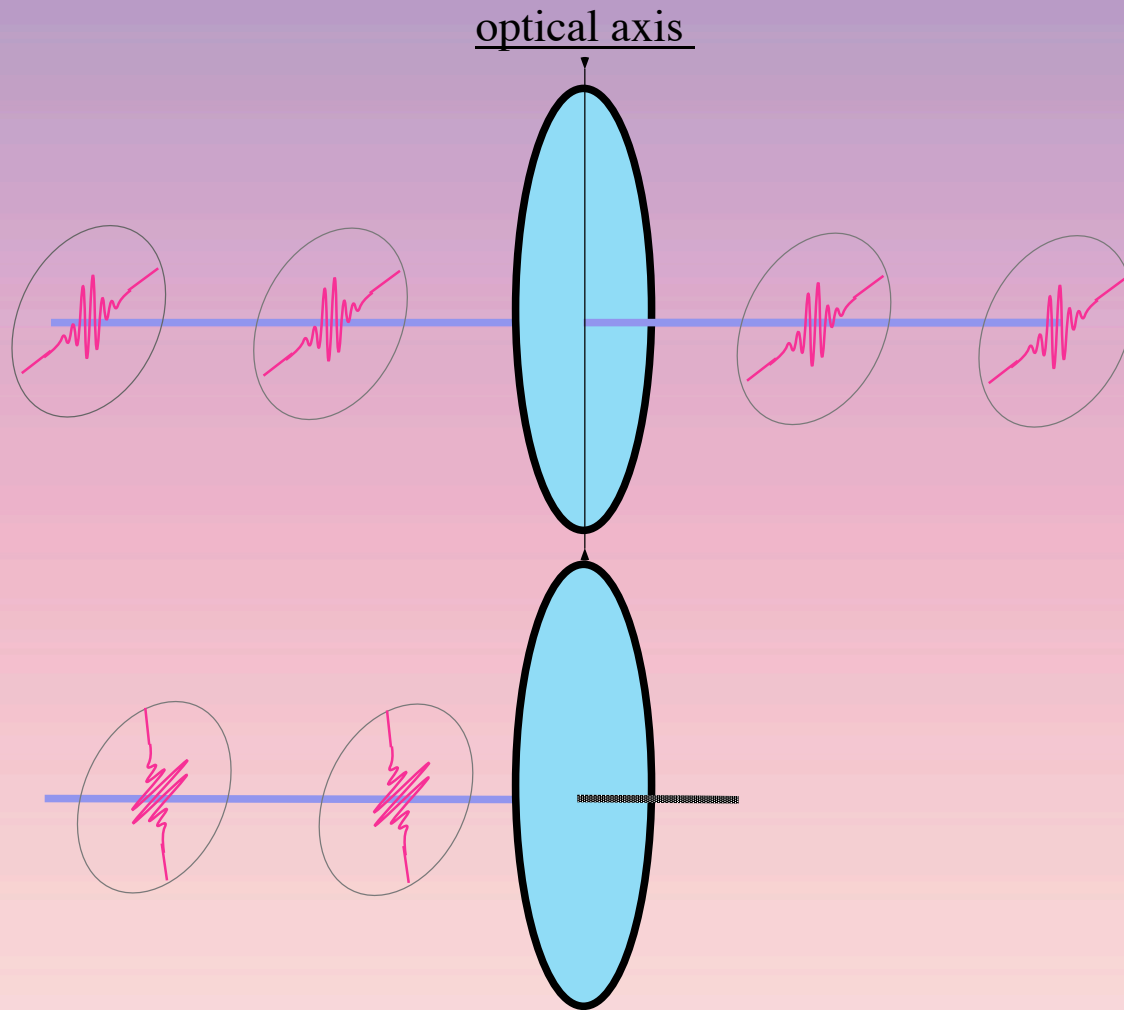
$|0\rangle \times |0\rangle = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$   $|0\rangle \times |1\rangle = \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix}$  ...  $|1\rangle \times |1\rangle = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$



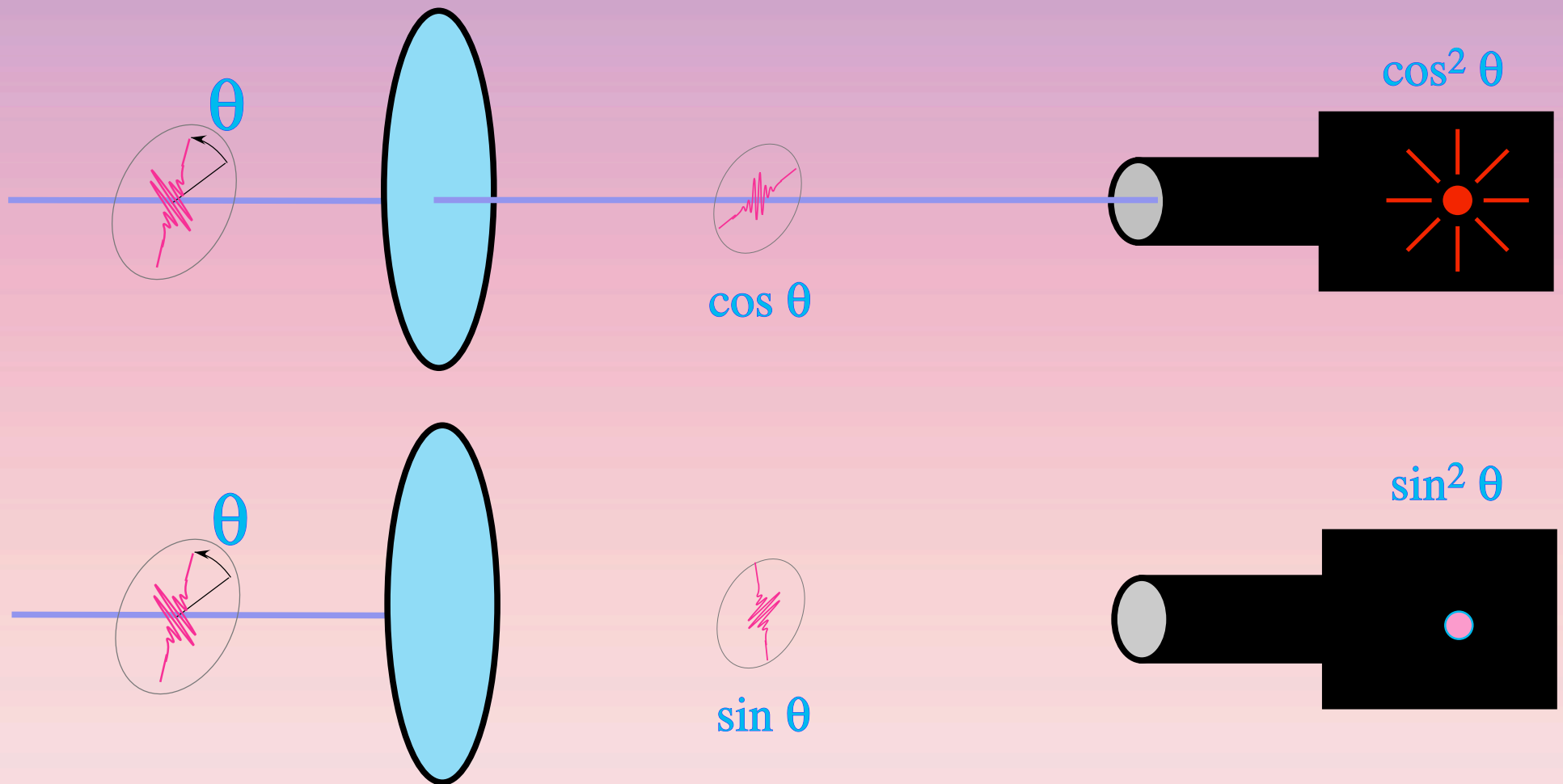
**(2)**

# **Quantum Measurements**

# Polarizing Filter



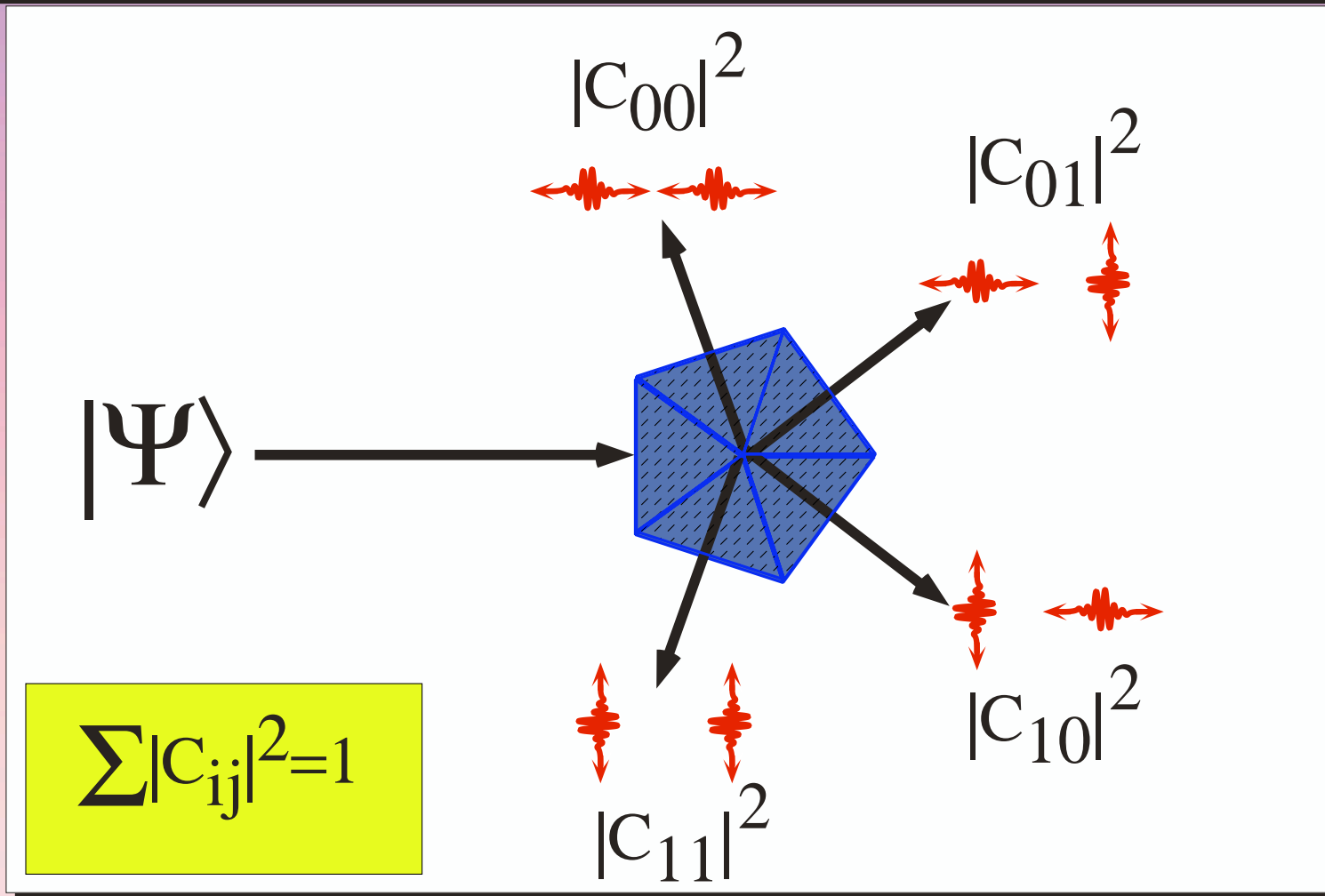
# Polarizing Filter and photodetectors



# Photons

# Projective Measurements

$$|\Psi\rangle = C_{00} \begin{array}{c} \leftarrow \text{---} \text{---} \text{---} \rightarrow \\ \leftarrow \text{---} \text{---} \text{---} \rightarrow \end{array} + C_{01} \begin{array}{c} \leftarrow \text{---} \text{---} \rightarrow \\ \updownarrow \end{array} + C_{10} \begin{array}{c} \updownarrow \\ \leftarrow \text{---} \text{---} \rightarrow \end{array} + C_{11} \begin{array}{c} \updownarrow \\ \updownarrow \end{array}$$



# standard notations

Projector:  $P_\psi = |\psi\rangle\langle\psi|$

$$P = \sum P_m, \text{ where } P_m P_{m'} = \delta_{m,m'} P_m$$

Measurement:  $\{ P_m \}$

$$\sum P_m = I, \quad P_m P_{m'} = \delta_{m,m'} P_m$$

Measuring:  $p(m) = \langle\phi|P_m|\phi\rangle$

Resulting:  $P_m|\phi\rangle/\sqrt{p(m)}$

## example

Projectors:  $P_0 = |0\rangle\langle 0| = \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}$

$$P_1 = |1\rangle\langle 1| = \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$$

Measurement:  $\{ P_0, P_1 \}$

$$\sum P_m = I, \quad P_0 P_1 = P_1 P_0 = 0$$

Measuring:  $p(m) = \langle \phi | P_m | \phi \rangle$

$$\alpha^* \langle 0 | + \beta^* \langle 1 | P_0 (\alpha | 0 \rangle + \beta | 1 \rangle) = \alpha^* \alpha = |\alpha|^2$$

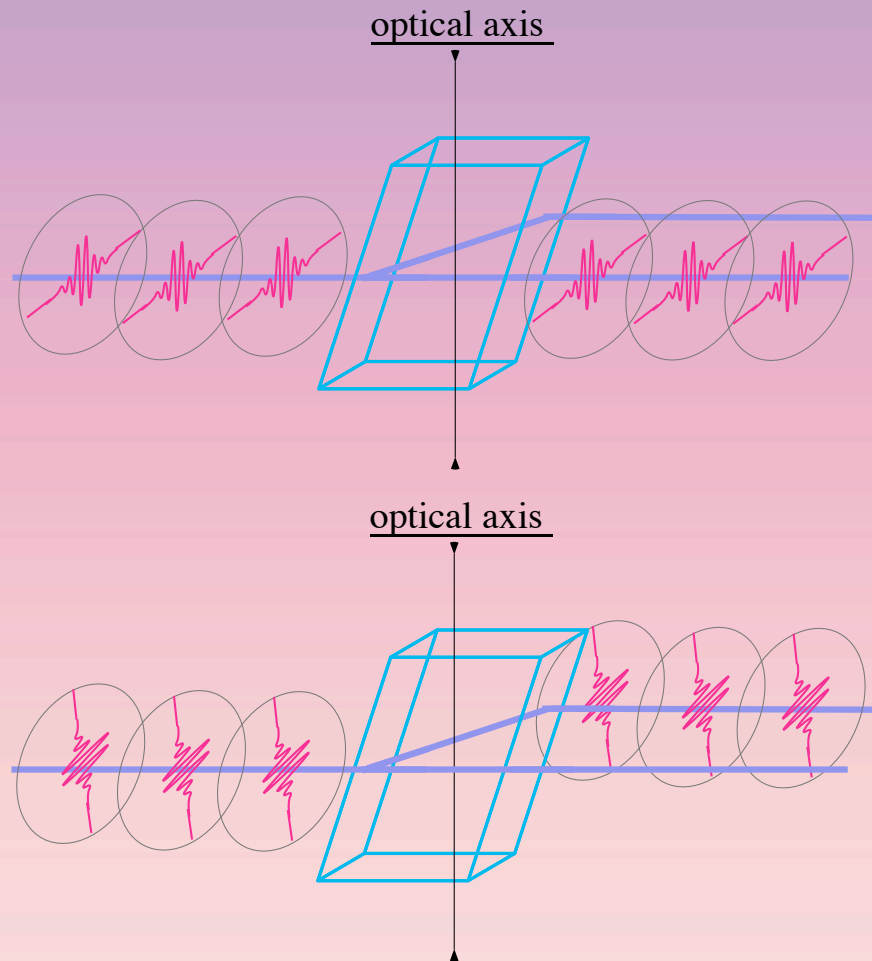
$$\alpha^* \langle 0 | + \beta^* \langle 1 | P_1 (\alpha | 0 \rangle + \beta | 1 \rangle) = \beta^* \beta = |\beta|^2$$

**(3)**

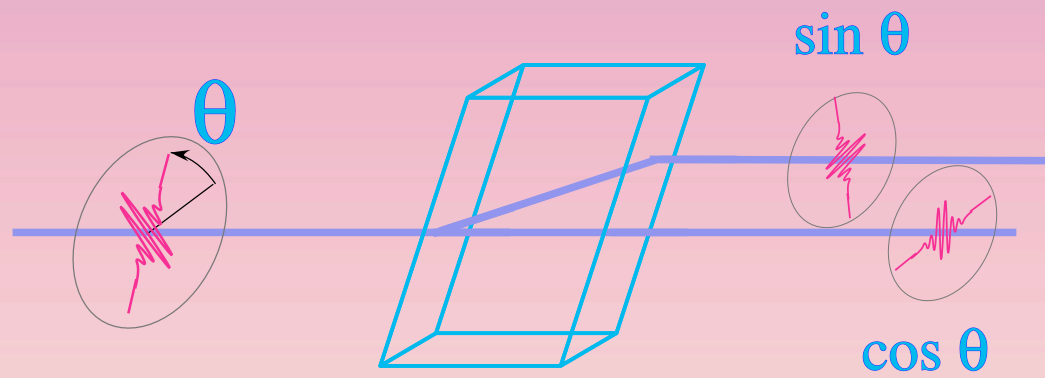
# **Quantum Computations**



# Calcite Crystal



# Calcite Crystal



# Photons

# Quantum Evolution: Unitary Operators

$$|\Psi\rangle \xrightarrow{\boxed{U}} |\Psi'\rangle$$

$$\text{↔} \xrightarrow{\boxed{U}} |\Psi_0\rangle$$

$$\text{↕} \xrightarrow{\boxed{U}} |\Psi_1\rangle$$

$$C_0 \text{↔} + C_1 \text{↕} \xrightarrow{\boxed{U}} C_0 |\Psi_0\rangle + C_1 |\Psi_1\rangle$$

## standard notations

Unitary Operators:  $U^\dagger U = U U^\dagger = I$

linear:  $U(\alpha|0\rangle + \beta|1\rangle) = \alpha U|0\rangle + \beta U|1\rangle$

preserves scalar products:

$$(\langle \psi | U^\dagger)(U | \phi \rangle)$$

$$= \langle \psi | U^\dagger U | \phi \rangle = \langle \psi | I | \phi \rangle = \langle \psi | \phi \rangle$$

## example: calcite crystal

Basis states:  $|\phi\pi\rangle$ ,  $\phi$ : polarizaton  $\{0,1\}$   
 $\pi$ : path  $\{\text{low,high}\}$

Input states:  $(\alpha|0\rangle + \beta|1\rangle)|\text{low}\rangle$

Operator:  $U: |0x\rangle \rightarrow |0x\rangle$       $U = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$

$|1x\rangle \rightarrow |1\neg x\rangle$

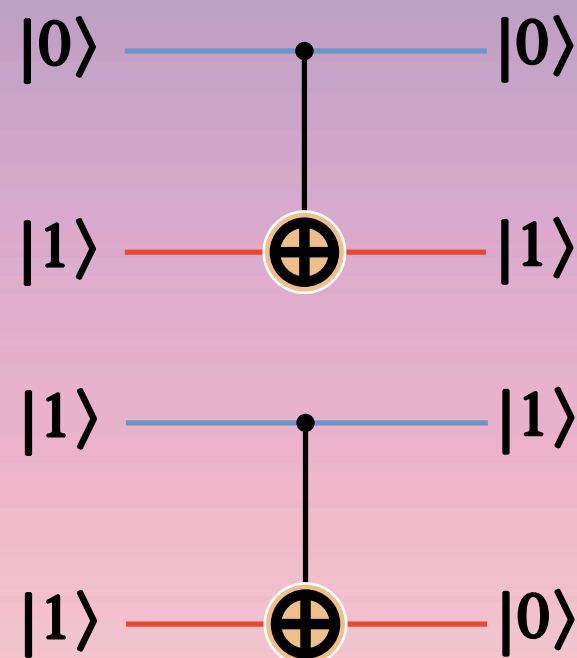
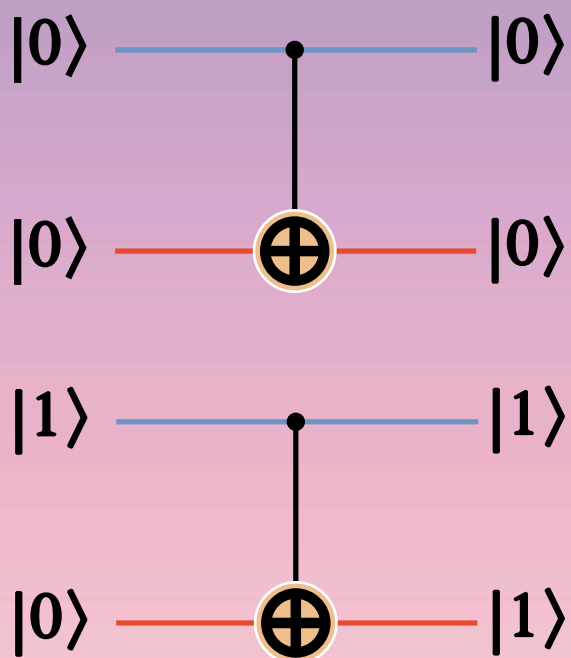
Output states:  $\alpha|0 \text{ low}\rangle + \beta|1 \text{ high}\rangle$

# example: Hadamard

$$\begin{array}{l} |0\rangle \text{ --- } \boxed{\text{H}} \text{ --- } (|0\rangle + |1\rangle)/\sqrt{2} \\ |1\rangle \text{ --- } \boxed{\text{H}} \text{ --- } (|0\rangle - |1\rangle)/\sqrt{2} \end{array}$$

$$\boxed{\text{H}} = 1/\sqrt{2} \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$$

# example: Control-NOT



A quantum circuit diagram of a Control-NOT gate (top blue line, bottom red line) is shown to be equivalent to the following matrix:

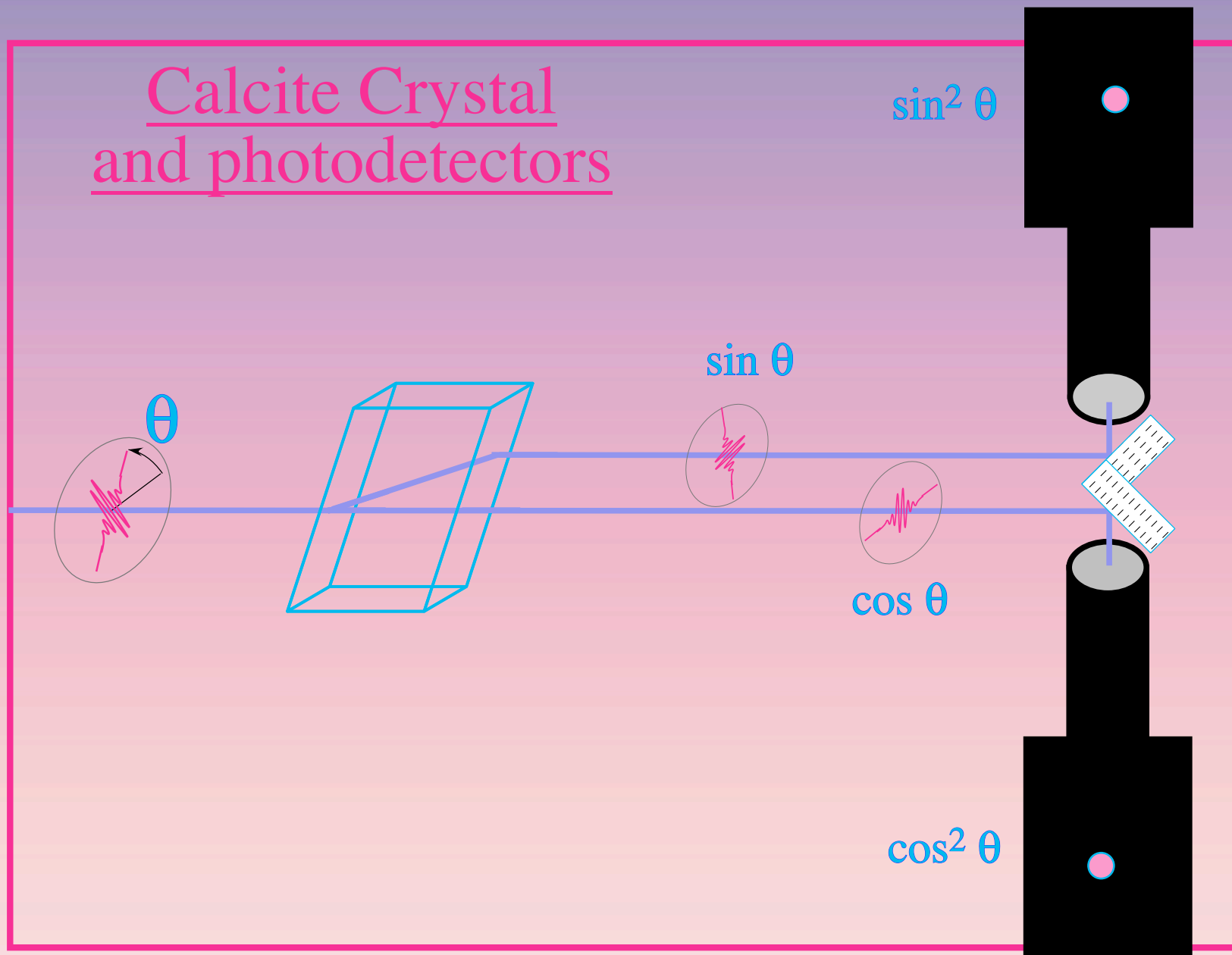
$$= \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$



**(4)**

# **General Measurements**

# Calcite Crystal and photodetectors



# standard notations

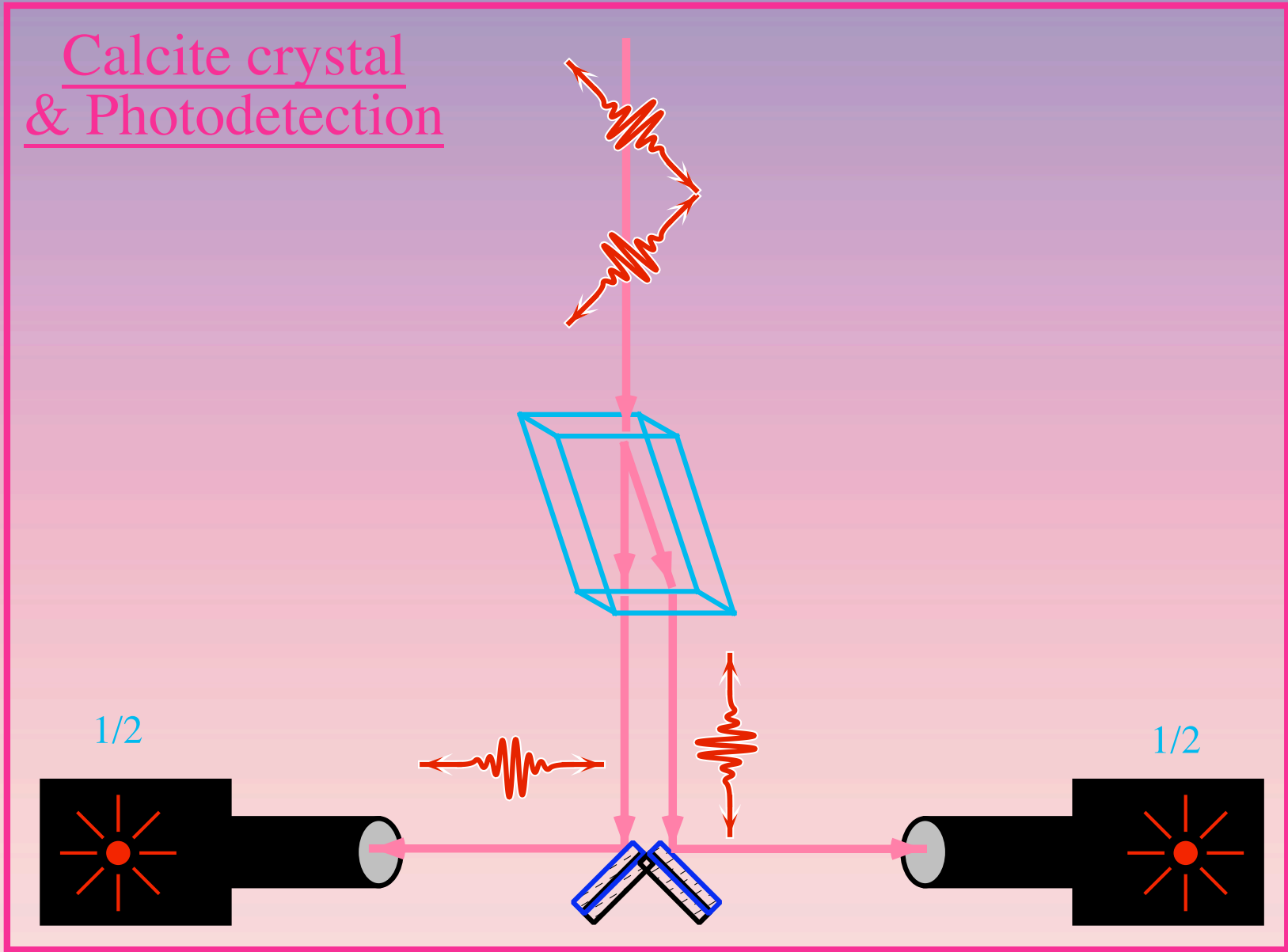
Measurement:  $\{ P_m U \}$

$$\sum P_m = I, \quad P_m P_{m'} = \delta_{m,m'} P_m$$

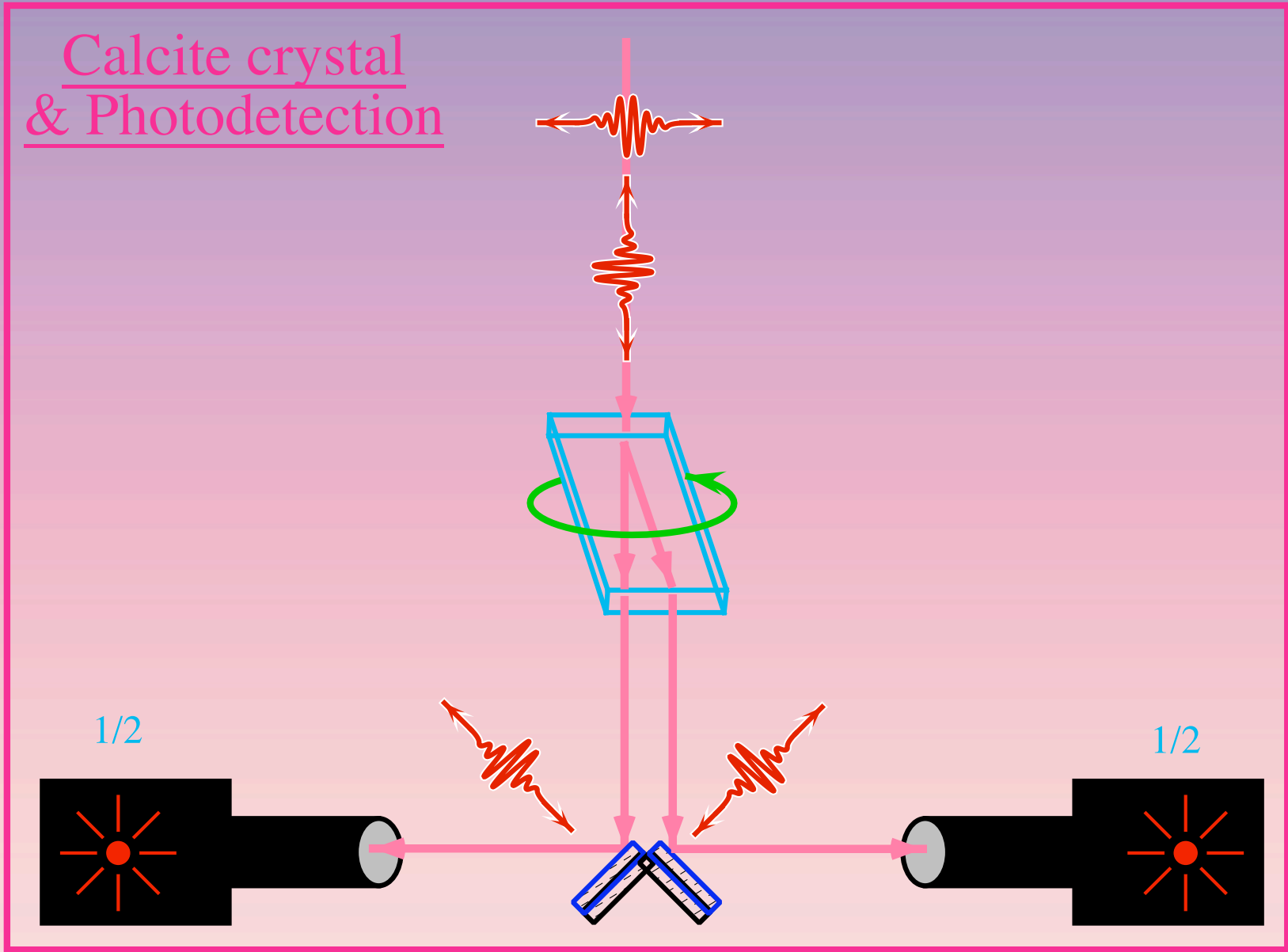
Measuring:  $p(m) = \langle \phi | \langle 0 | U^\dagger P_m U | \phi \rangle | 0 \rangle$

Resulting:  $P_m U | \phi \rangle | 0 \rangle / \sqrt{p(m)}$

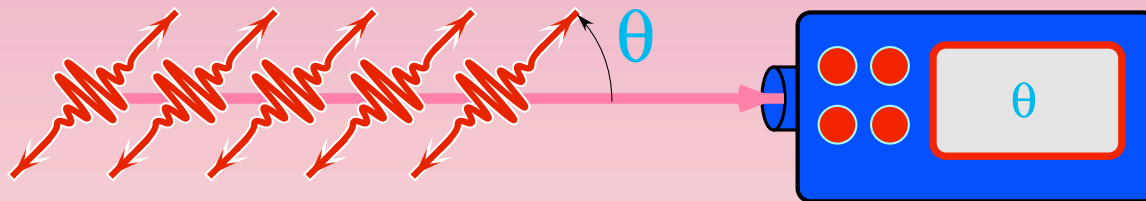
# Calcite crystal & Photodetection



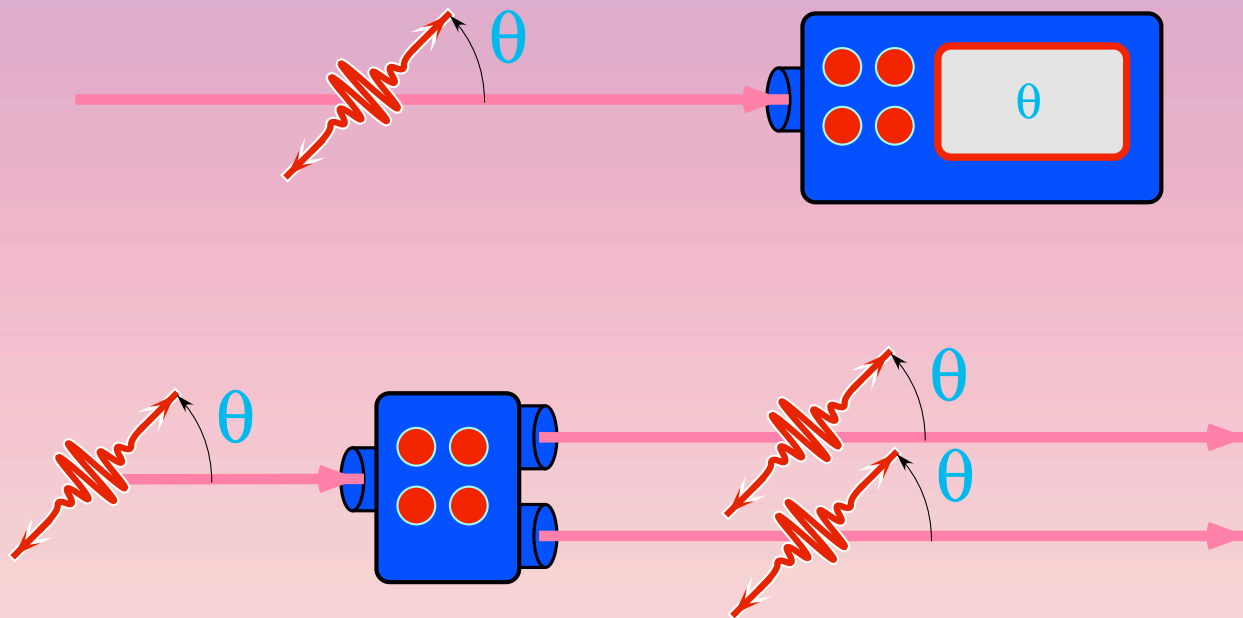
# Calcite crystal & Photodetection



POSSIBLE!



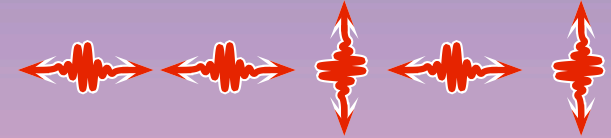
# IMPOSSIBLE!



# Classical & Quantum Information

00110111000110 Classical

Quantum



Copying:

Yes

NO

Measuring:

Yes

partial

Broadcasting:

Yes

NO

Superposing:

NO

Yes

Interfering:

NO

Yes



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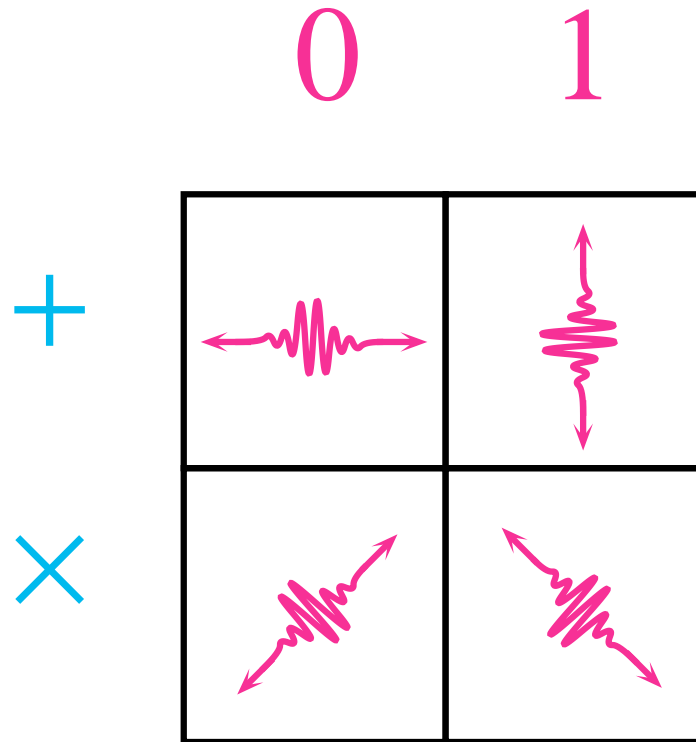


**(8)**

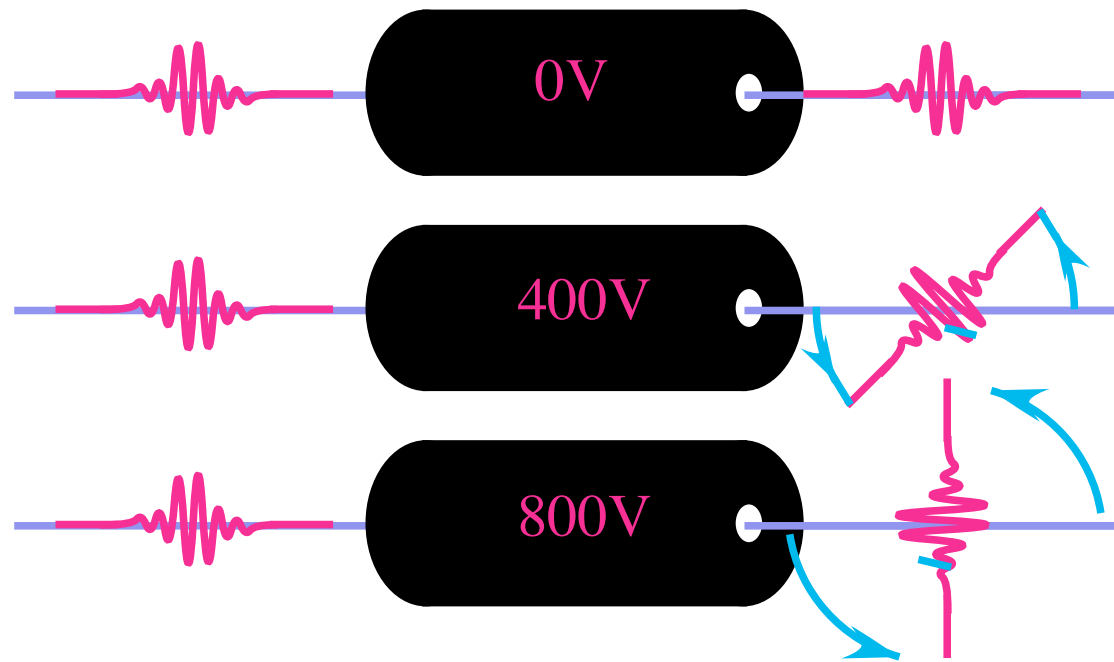
# **Quantum Cryptography**

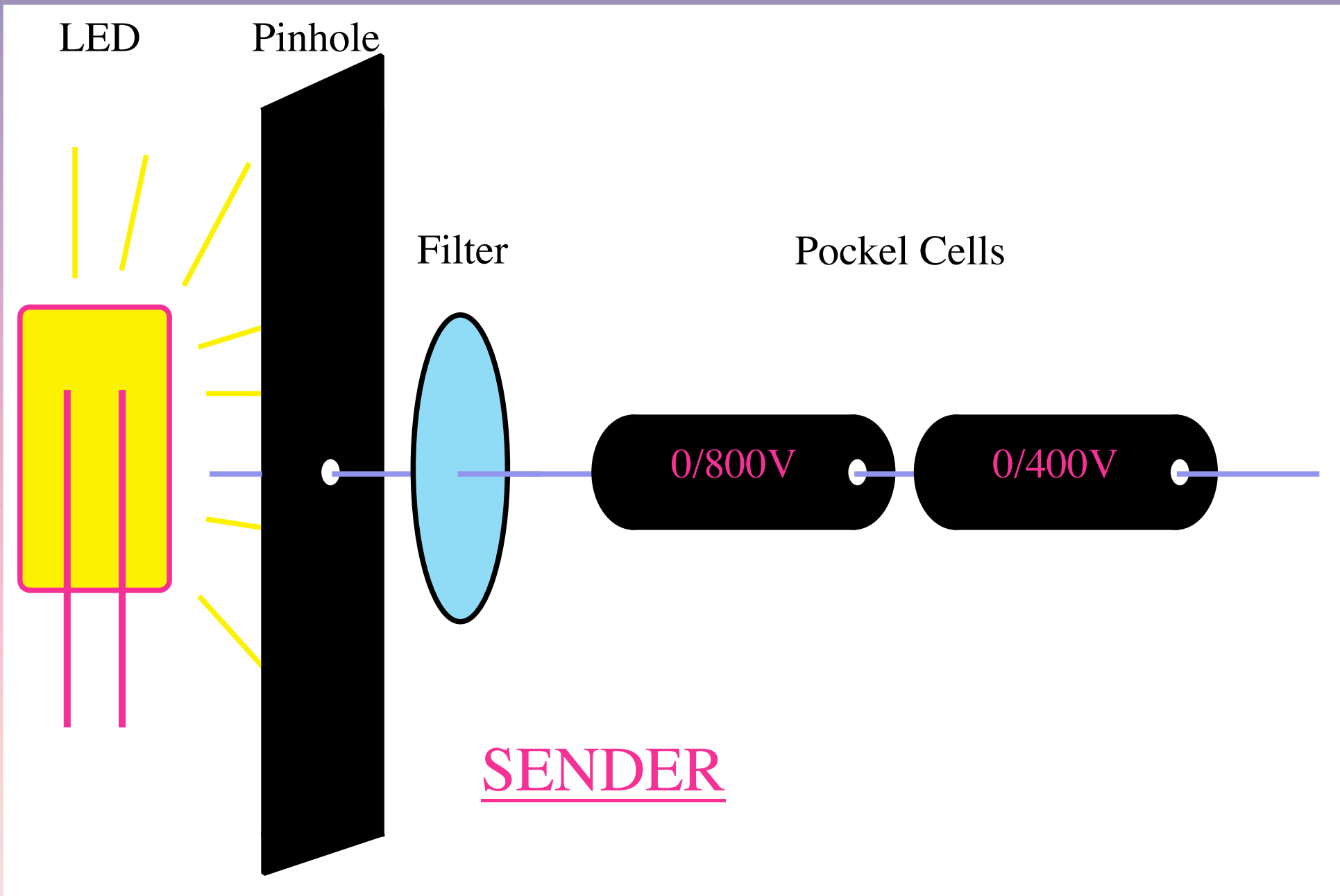
## Key distribution

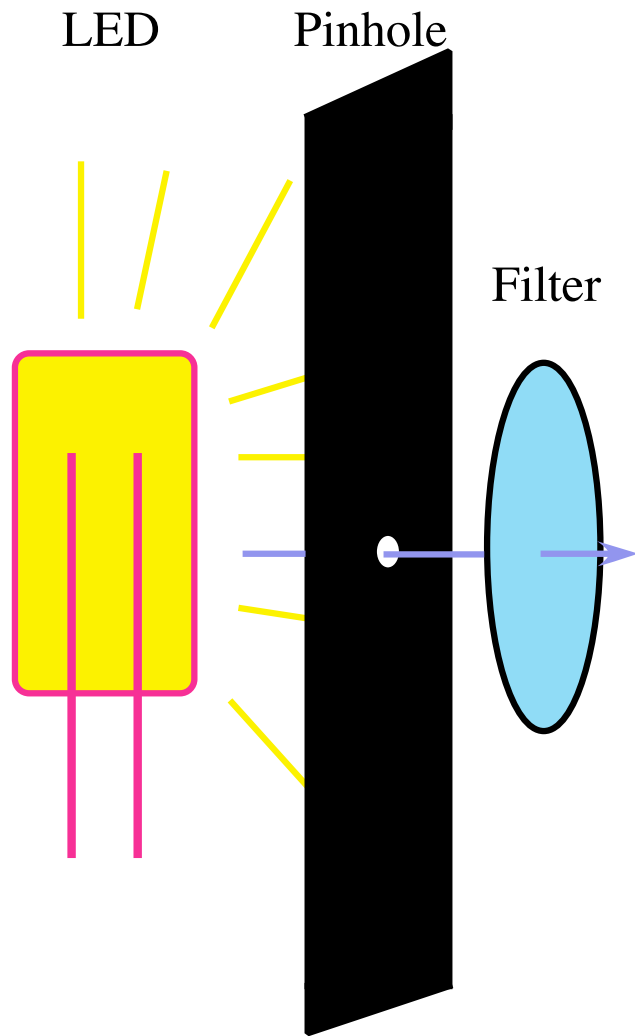
# Ambiguous Coding Scheme



# Pockel Cells







### Light source:

$\sim 1/10$  photon per pulse

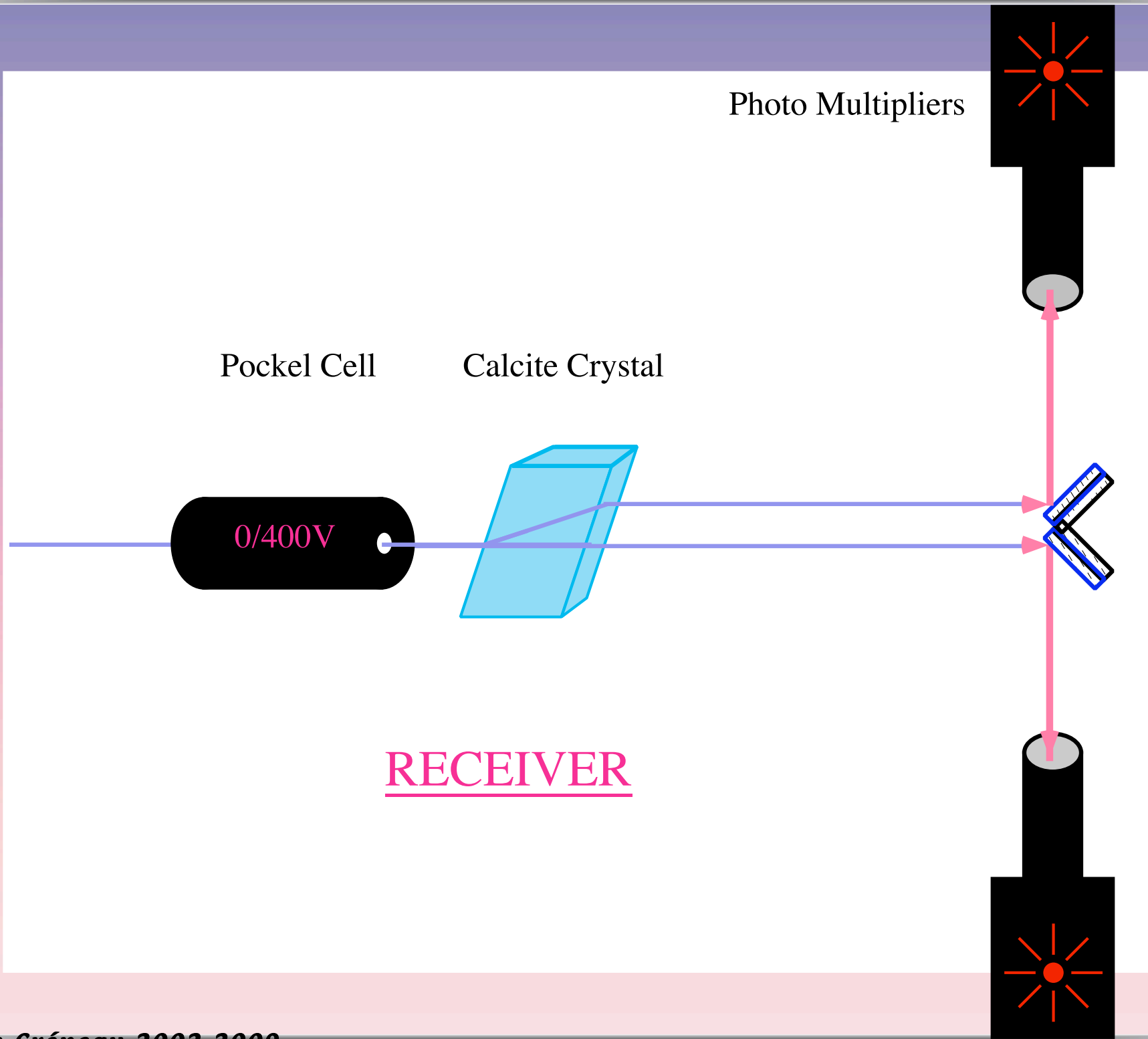
$n = \#$ photons per pulse follows a

Poisson distribution  $\Pr(n \leq x) = 1 - e^{-x/10}$

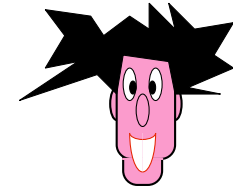
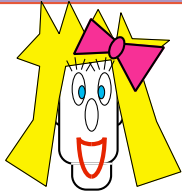
### Problem:

- may transmit multiple correlated

polarized photons



# Q-distribution of keys



A: 0 1 1 0 0 1 0 0 1 1 0 1 0 0 0 1 1 1 0 1 1 0 0 0  
 × + × + + + × × × × + + + + × × × + × + + + × +

B: × × + + × + + + × + + × × × + × × × + + × + × +  
 0 0 1 0 0 1 0 0 1 0 0 0 0 1 1 1 0 0 0 1 1 0 0 0

A: × + × + + + × × × × + + + + × × × + × + + + × +

B: 0   0  1   1  0     1 0   1  0 0 0

B: 0 0 1 1 0 1 0 1 0 0 0

A: 0 0 1 1 0 1 1 1 1 0 0 0

A: 0 1 0 1 0

B: = = = ≠ =

B: 0 1 1 1 0 0

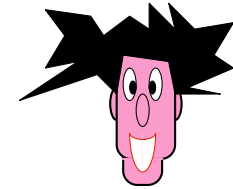
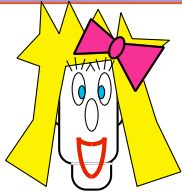
A: 0 1 1 1 0 0

20%

## Bennett- Brassard

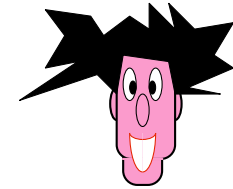
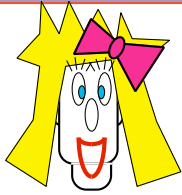


# Q-distribution of keys



A: 0 1 1 0 0 1 0 0 1 1 0 1 0 0 0 1 1 1 0 1 1 0 0 0  
× + × + + + × × × × + + + + × × × + × + + + × +

# Q-distribution of keys

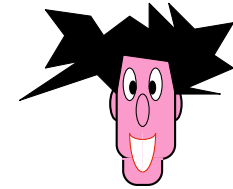
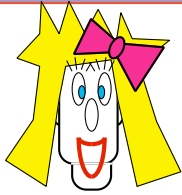


**A:** 0 1 1 0 0 1 0 0 1 1 0 1 0 0 0 1 1 1 0 1 1 0 0 0  
× + × + + + × × × × + + + + × × × + × + + + × +

**B:** × × + + × + + + × + + × × × + × × × + + × + × +  
0 0 1 0 0 1 0 0 1 0 0 0 0 1 1 1 0 0 0 1 1 0 0 0

---

# Q-distribution of keys



**A:** 0 1 1 0 0 1 0 0 1 1 0 1 0 0 0 1 1 1 0 1 1 0 0 0

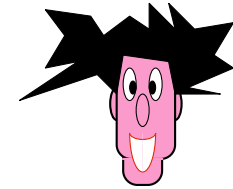
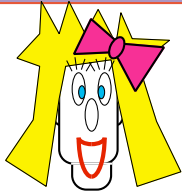
× + × + + + × × × × + + + + × × × + × + + + × +

**B:** × × + + × + + + × + + × × × + × × × + + × + × +

0 0 1 0 0 1 0 0 1 0 0 0 0 1 1 1 0 0 0 1 1 0 0 0

**A:** × + × + + + × × × × + + + + × × × + × + + + × +

# Q-distribution of keys



**A:** 0 1 1 0 0 1 0 0 1 1 0 1 0 0 0 1 1 1 0 1 1 0 0 0  
 × + × + + + × × × × + + + + × × × + × + + + × +

**B:** × × + + × + + + × + + × × × + × × × + + × + × +  
 0 0 1 0 0 1 0 0 1 0 0 0 0 1 1 1 0 0 0 1 1 0 0 0

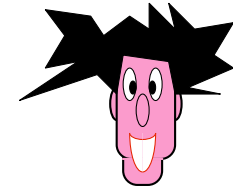
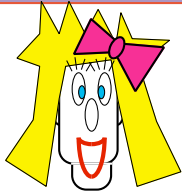
**A:** × + × + + + × × × × + + + + × × × + × + + + × +

**B:** 0   0  1   1  0     1 0   1  0 0 0

**B:** 0 0 1 1 0 1 0 1 0 0 0

**A:** 0 0 1 1 0 1 1 1 1 0 0 0

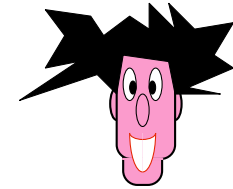
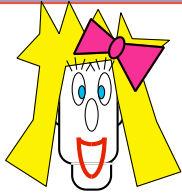
# Q-distribution of keys



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<b>B:</b>	0	0	1	1	0	1	0	1	0	0	0
<b>A:</b>	0	0	1	1	0	1	1	1	0	0	0

# Q-distribution of keys



B: 0 0 1 1 0 1 0 1 0 0 0

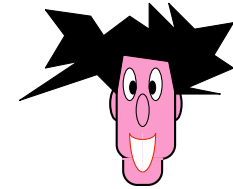
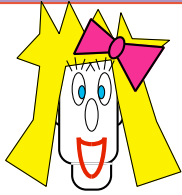
A: 0 0 1 1 0 1 1 1 0 0 0

A: 0 1 0 1 0

B: = = = ↑ =

20%

# Q-distribution of keys



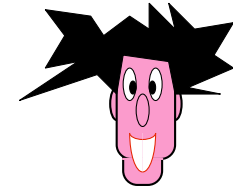
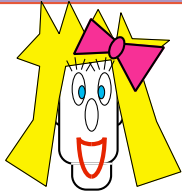
B: = = = ↑ =

B: 0 1 1 1 0 0

A: 0 1 1 1 0 0

20%

# Q-distribution of keys



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<b>B:</b>	0	1	1	1	0 0
<b>A:</b>	0	1	1	1	0 0

20%



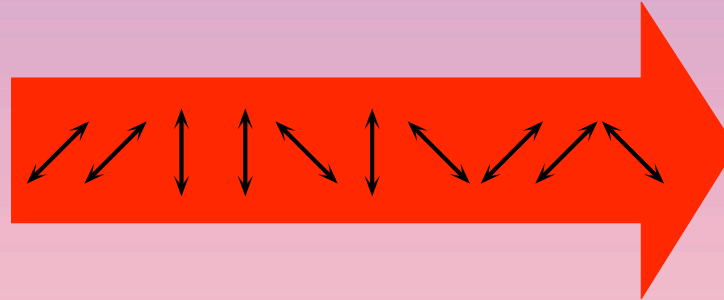
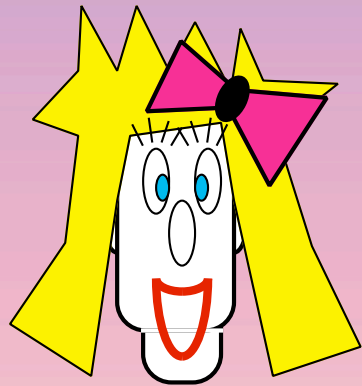
# Q-distribution of keys

.....

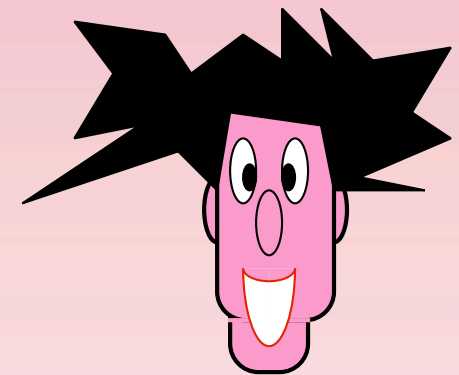
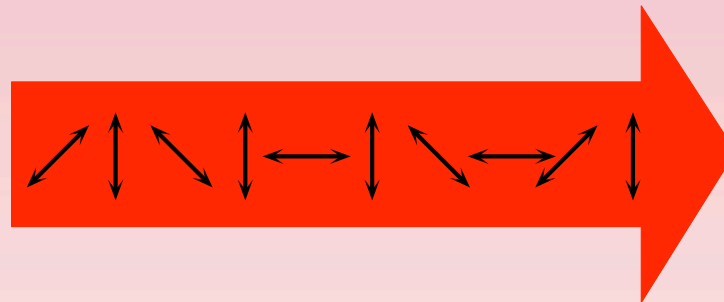
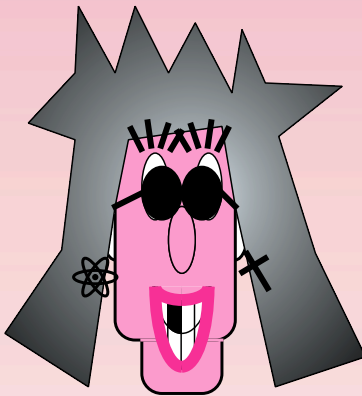
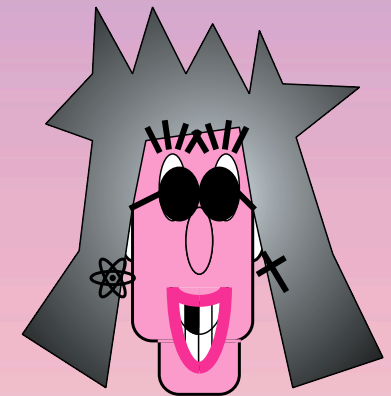
- Produces raw classical key
- Observed error rate indicates amount of eavesdropper information
- Error-correction is used to fix errors
- Random hash function is used to distill a smaller secret classical key

.....

# Information <--> Errors (intercept resend)

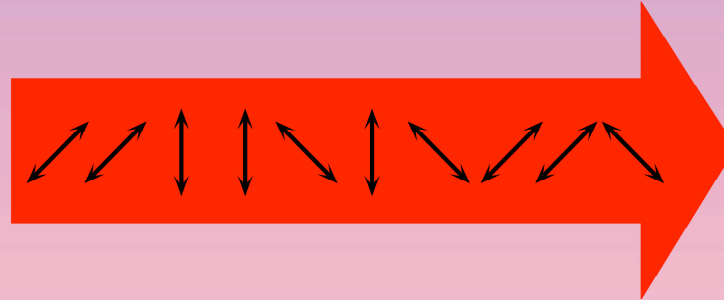
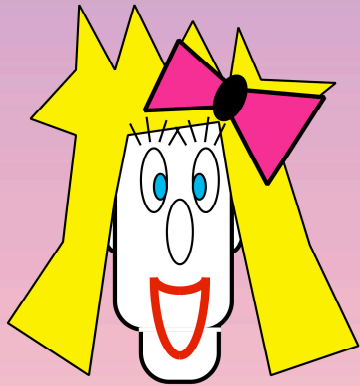


x + x + + + x + x +

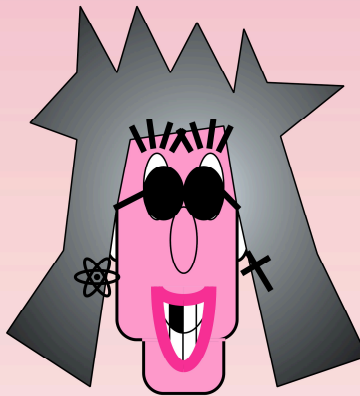


x + x + + + x + x +

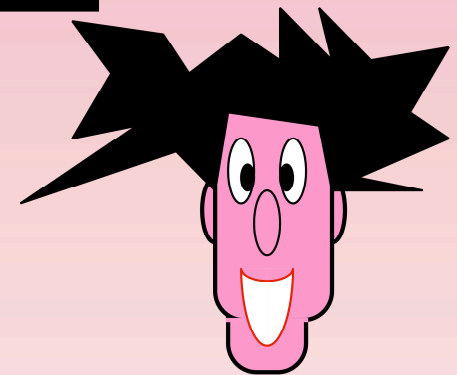
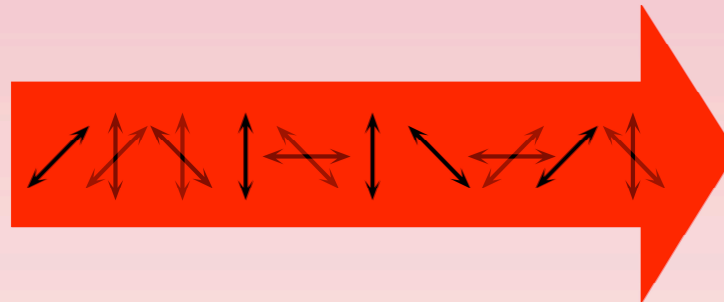
# Information <--> Errors (general)



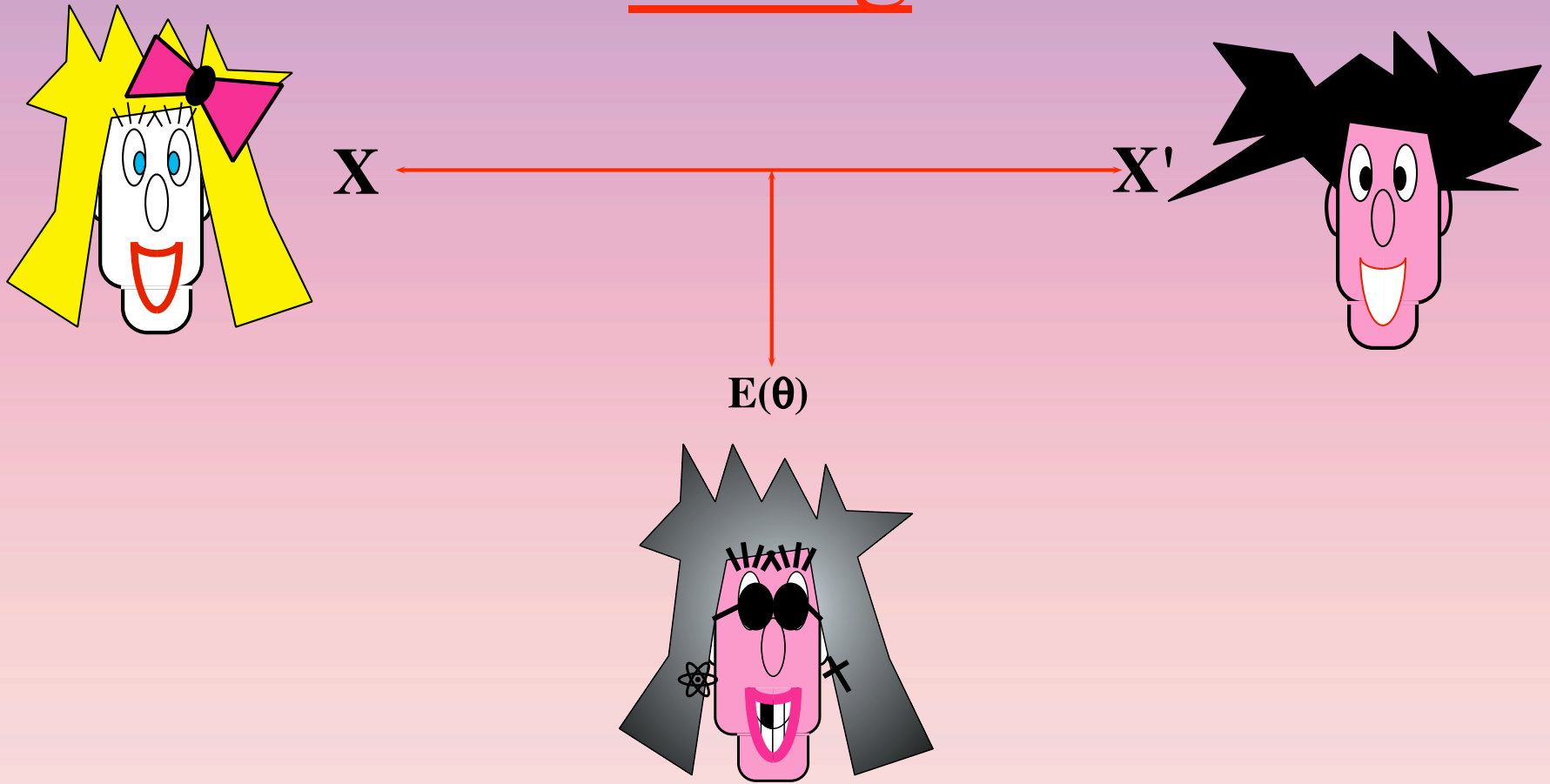
A  
r  
b  
i  
t  
r  
a  
r  
y  
M  
e  
a  
s  
u  
r  
e  
m  
e  
n  
t



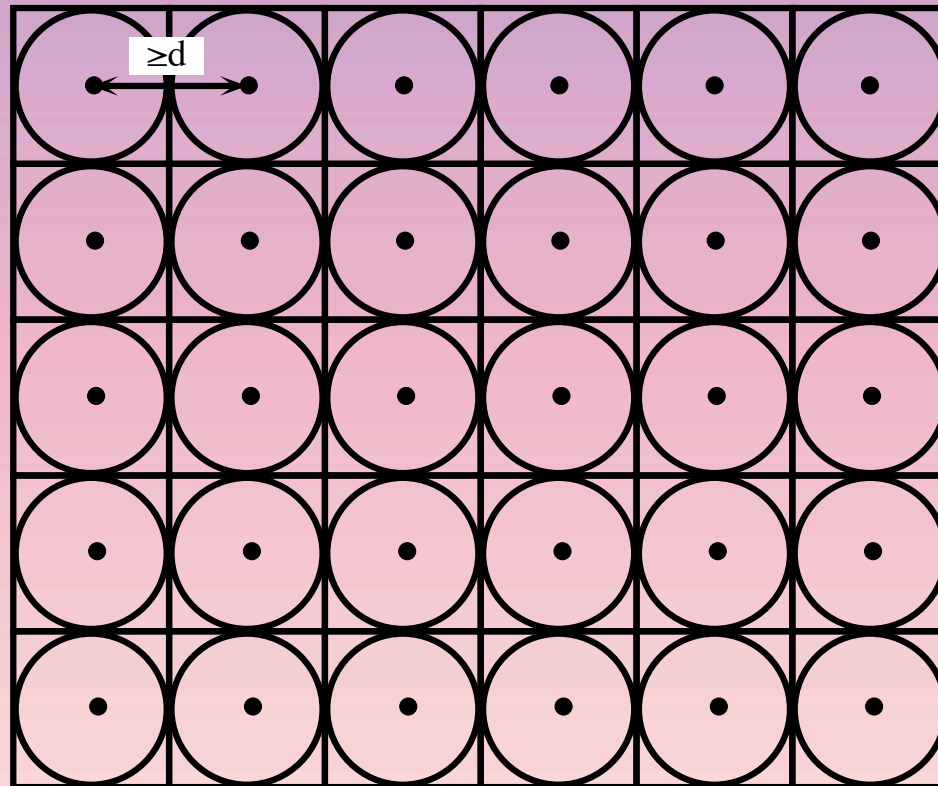
L  
e  
f  
t  
o  
v  
e  
r



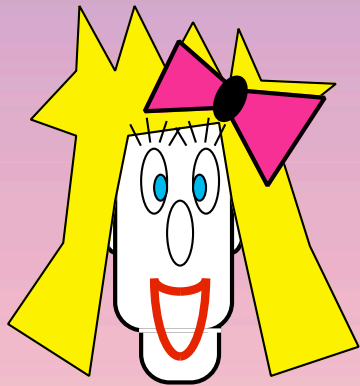
# Mostly Identical Partly Secret String



# (classical) error-correcting codes



# Identical Partly Secret String



**X**

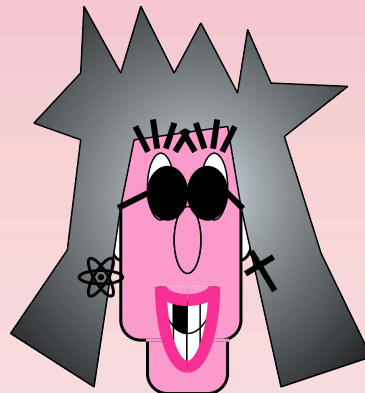


$$W := C \oplus X$$

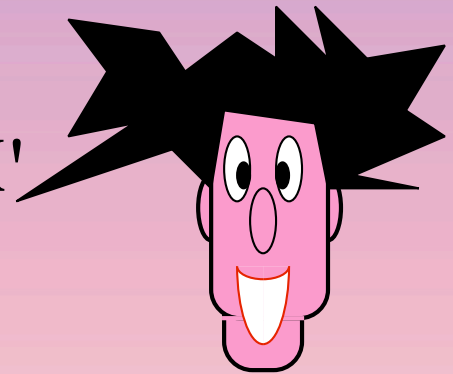
**W**



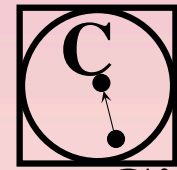
$$E := E(\theta) + W$$



**X'**



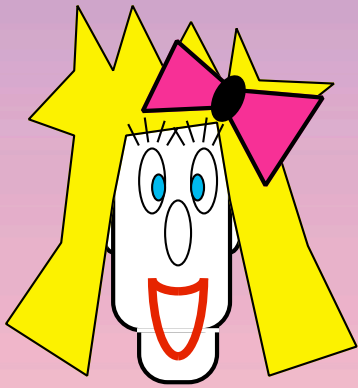
$$C' := W \oplus X'$$



**C'**

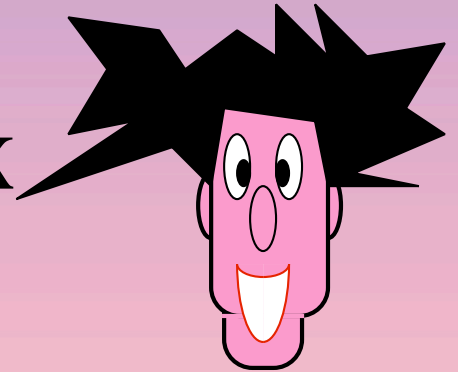
$$X := C \oplus W$$

# Identical Partly Secret String

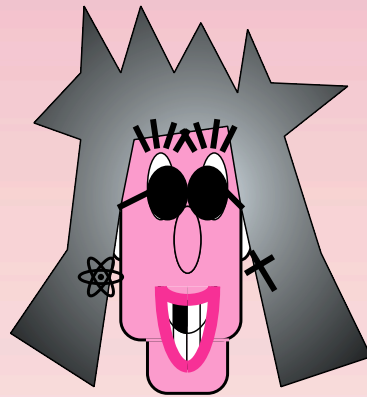


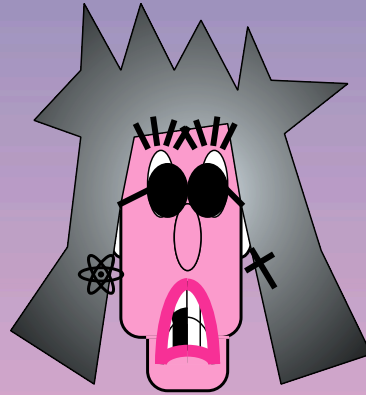
**X**

**X**



**E**



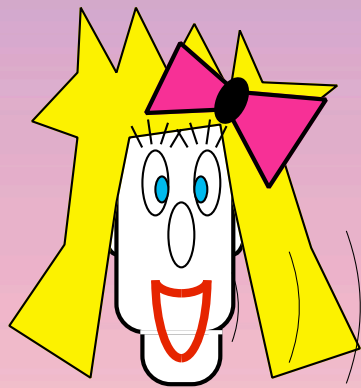


**in BB84, Eve's information at this point (before privacy amplification) contains:**

- **results of eavesdropping**
- **multi-photon pulses**
- **error correction**

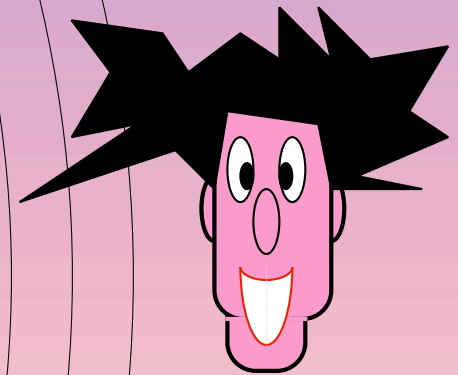


# Identical Secret Shorter String through Privacy Amplification

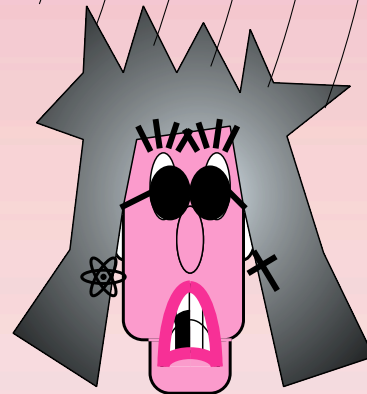


$$X \rightarrow h(X)$$

h



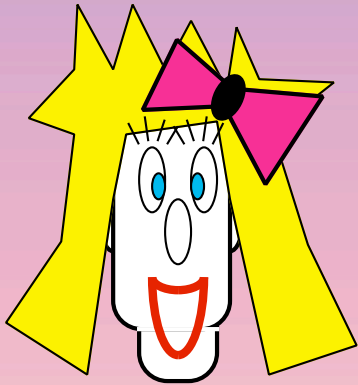
$$X \rightarrow h(X)$$



$$E \rightarrow_{h(E)}$$

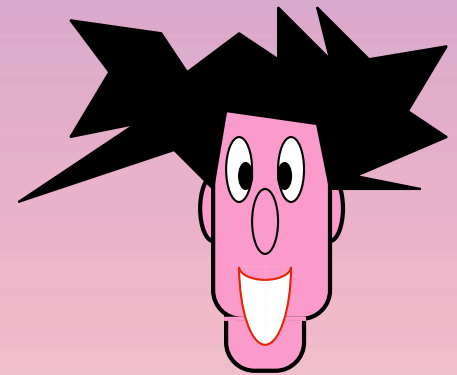
# BBCM

$$H( G(X) \mid E=e, G ) > |G(X)| - 2^{(|G(X)|-R(X|E=e))}$$

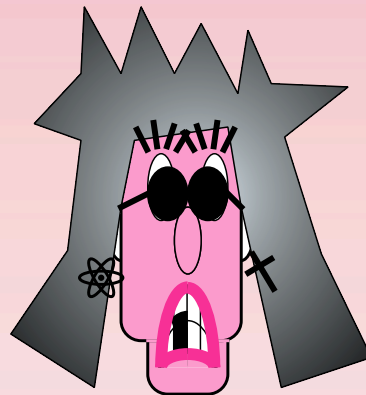


$k:=g(X)$

g



$k:=g(X)$



??????????  
 $k':=g(E)$   
??????????