

COMP-647B Advanced Cryptography

Problem set #1

due on Tuesday January 30, 2007

1. MA vs AM

Prove that $\mathbf{MA} \subseteq \mathbf{AM}$.

2. Code equivalence in ZK.

We say that two binary matrices G, G' generate two *equivalent* linear codes $C = \text{span}(G)$ and $C' = \text{span}(G')$ if there exists a permutation matrix P (each row and column contain all “0” except for one “1”) and a base change matrix S (full rank) such that

$$G' = SGP.$$

Give a *perfect* zero-knowledge interactive proof for the language

$$L_{\text{eqv}} = \{ (G, G') \mid G, G' \text{ generate } \textit{equivalent} \text{ linear codes} \}.$$

3. Hamiltonian in ZK.

Give a *computational* zero-knowledge interactive proof for the Hamiltonian circuit problem under suitable computational assumption. (A directed graph G is Hamiltonian if its edges contain a circuit visiting each vertex exactly once.)

4. RSA integers in ZK.

4.A Let **RSA** be integers with exactly two distinct prime factors. Give a Zero-Knowledge interactive proof for the **RSA** numbers.

HINT:

We define **RSA** = $\{n \mid n=pq \text{ where } p,q \text{ are distinct primes}\}$.

You may use without proof the following two results:

Theorem 1. If $n=p_1^{\alpha_1}p_2^{\alpha_2}\dots p_k^{\alpha_k}$ (which is not a square) has exactly k distinct prime factors then exactly 2^{1-k} of the x in Z_n^* with Jacobi symbol $+1$ are quadratic residues mod n .

Theorem 2. Let n be a composite number. If $n=p_1p_2\dots p_k$ is a product of distinct primes then every x in Z_n^* has an n^{th} root mod n , i.e. a y such that $x \equiv y^n \pmod n$. On the contrary, if $p_i=p_j$ for some $1 \leq i < j \leq k$, then at most half the x in Z_n^* has an n^{th} root mod n .

Construct two zero-knowledge proofs for the languages **WRSA** of **Weak-RSA** numbers and **SF** of square-free numbers:

$$\mathbf{WRSA} = \{n \mid n=p^\alpha q^\beta \text{ where } p,q \text{ are distinct primes and } \alpha,\beta > 0\}$$

$$\mathbf{SF} = \{n \mid n=p_1p_2\dots p_k \text{ is a product of distinct primes}\}$$

Notice that **RSA** = **WRSA** \cap **SF**.

4.B Finally, if we define

$$\mathbf{BLUM} = \{n \mid n=pq \text{ where } p \equiv q \equiv 3 \pmod 4, \text{ are distinct primes}\}$$

then prove that **BLUM** has a *statistical ZK* interactive proof.

HINT:

Define the **Weak-BLUM** integers as

$$\mathbf{WBLUM} = \{n \mid -1 \text{ is in } \text{QNR}_n[+1] \}$$

Notice that **BLUM** = **RSA** \cap **WBLUM**.