Chapter 4
Greedy Algorithms
4.5 Minimum Spanning Tree
Minimum spanning tree. Given a connected graph $G = (V, E)$ with real-valued edge weights $c_e$, an MST is a subset of the edges $T \subseteq E$ such that $T$ is a spanning tree whose sum of edge weights is minimized.

\[ G = (V, E) \quad T, \quad \sum_{e \in T} c_e = 50 \]

Cayley's Theorem. There are $n^{n-2}$ spanning trees of $K_n$. can't solve by brute force
Applications

MST is fundamental problem with diverse applications.

- Network design.
  - telephone, electrical, hydraulic, TV cable, computer, road

- Approximation algorithms for NP-hard problems.
  - traveling salesperson problem, Steiner tree

- Indirect applications.
  - max bottleneck paths
  - LDPC codes for error correction
  - image registration with Renyi entropy
  - learning salient features for real-time face verification
  - reducing data storage in sequencing amino acids in a protein
  - model locality of particle interactions in turbulent fluid flows
  - autoconfig protocol for Ethernet bridging to avoid cycles in a network

- Cluster analysis.
Greedy Algorithms

**Kruskal's algorithm.** Start with $T = \emptyset$. Consider edges in ascending order of cost. Insert edge $e$ in $T$ unless doing so would create a cycle.

**Reverse-Delete algorithm.** Start with $T = E$. Consider edges in descending order of cost. Delete edge $e$ from $T$ unless doing so would disconnect $T$.

**Prim's algorithm.** Start with some root node $s$ and greedily grow a tree $T$ from $s$ outward. At each step, add the cheapest edge $e$ to $T$ that has exactly one endpoint in $T$.

**Remark.** All three algorithms produce an MST.
Greedy Algorithms

Simplifying assumption. All edge costs $c_e$ are **distinct**.

Cut property. Let $S$ be any subset of nodes, and let $e$ be the min cost edge with exactly one endpoint in $S$. Then the MST contains $e$.

Cycle property. Let $C$ be any cycle, and let $f$ be the max cost edge belonging to $C$. Then the MST does not contain $f$.

- $e$ is in the MST
- $f$ is not in the MST
**Cycles and Cuts**

**Cycle.** Set of edges the form a-b, b-c, c-d, ..., y-z, z-a.

\[
\text{Cycle } C = 1-2, 2-3, 3-4, 4-5, 5-6, 6-1
\]

**Cut-set.** A cut is a subset of nodes S. The corresponding cut-set D is the subset of edges with exactly one endpoint in S.

\[
\text{Cut } S = \{ 4, 5, 8 \}, \quad \text{Cut-set } D = 5-6, 5-7, 3-4, 3-5, 7-8
\]
Claim. A cycle and a cut-set intersect in an even number of edges.

Pf. (by picture)
**Simplifying assumption.** All edge costs $c_e$ are **distinct**.

**Cut property.** Let $S$ be any subset of nodes, and let $e$ be the min cost edge with exactly one endpoint in $S$. Then any MST $T^*$ must contain $e$.

**Pf. (exchange argument)**
- Suppose $e$ does not belong to MST $T^*$, and let's see what happens.
- Adding $e$ to $T^*$ creates a cycle $C$ in $T^*$.
- Edge $e$ is both in the cycle $C$ and in the cut-set $D$ corresponding to $S$ $\Rightarrow$ there exists another edge, say $f$, that is in both $C$ and $D$.
- $T' = T^* \cup \{e\} - \{f\}$ is also a spanning tree.
- Since $c_e < c_f$, cost($T'$) $<$ cost($T^*$).
- This is a contradiction. ▪
Simplifying assumption. All edge costs $c_e$ are distinct.

Cycle property. Let $C$ be any cycle in $G$, and let $f$ be the max cost edge belonging to $C$. Then any MST $T^*$ does not contain $f$.

Pf. (exchange argument)
- Suppose $f$ belongs to MST $T^*$, and let's see what happens.
- Deleting $f$ from $T^*$ creates a cut $S$ in $T^*$.
- Edge $f$ is both in the cycle $C$ and in the cut-set $D$ corresponding to $S$ \[ \Rightarrow \text{there exists another edge, say } e, \text{ that is in both } C \text{ and } D. \]
- $T' = T^* \cup \{e\} - \{f\}$ is also a spanning tree.
- Since $c_e < c_f$, cost($T'$) < cost($T^*$).
- This is a contradiction. \[ \blacksquare \]
Prim's Algorithm: Proof of Correctness

Prim's algorithm.  [Jarník 1930, Dijkstra 1957, Prim 1959]

- Initialize $S = \text{any node}$.
- Apply cut property to $S$.
- Add to $T$ the min cost edge in cut-set corresponding to $S$, and add one new explored node $u$ to $S$. 

![Diagram of Prim's algorithm](image-url)
Implementation: Prim's Algorithm

Implementation. Use a priority queue à la Dijkstra.

- Maintain set of explored nodes $S$.
- For each unexplored node $v$, maintain attachment cost $a[v] =$ cost of cheapest edge from $v$ to a node in $S$, and set $e[v] =$ cheapest edge.

```plaintext
Prim(G, c) {
    foreach (v ∈ V) a[v] ← ∞, e[v] ← Ø
    {Initialize an empty priority queue Q}
    foreach (v ∈ V) insert v onto Q
    S ← Ø, T ← Ø

    while (Q is not empty) {
        u ← extract min element from Q
        S ← S ∪ { u }, T ← T ∪ e[u]
        foreach (edge e = (u, v))
            if ((v ∉ S) and (c_e < a[v]))
                decrease priority a[v] to c_e in Q
                set e[v] to {e}
    }
}
Implementation. Use a priority queue à la Dijkstra.
- Maintain set of explored nodes $S$.
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<table>
<thead>
<tr>
<th>Minimum edge weight data structure</th>
<th>Time complexity (total)</th>
</tr>
</thead>
<tbody>
<tr>
<td>adjacency matrix, searching</td>
<td>$O(n^2)$</td>
</tr>
<tr>
<td>binary heap and adjacency list</td>
<td>$O((n+m) \log n) = O(m \log n)$</td>
</tr>
<tr>
<td>Fibonacci heap and adjacency list</td>
<td>$O(m + n \log n)$</td>
</tr>
</tbody>
</table>
Kruskal's Algorithm: Proof of Correctness

Kruskal's algorithm. [Kruskal, 1956]
Consider edges in ascending order of weight.
- **Case 1:** If adding $e$ to $T$ creates a cycle,
  discard $e$ according to cycle property.
- **Case 2:** Otherwise, insert $e = (u, v)$ into $T$ according to cut property where $S =$ set of nodes in $u$'s connected component.
Implementation: Kruskal's Algorithm

**Implementation.** Use the union-find data structure.
- Build set $T$ of edges in the MST.
- Maintain set for each connected component.
- $O(m \log n)$ for sorting and $O(m \alpha (m, n))$ for union-find.

\[
\begin{align*}
\text{essentially a constant}
\end{align*}
\]

\[m \leq n^2 \Rightarrow \log m \in O(\log n)\]

---

**Kruskal(G, c)** {

*Sort edges weights so that $c_1 \leq c_2 \leq \ldots \leq c_m$.*

$T \leftarrow \emptyset$

*foreach* $(u \in V)$ make a set containing singleton $u$

*for* $i = 1$ to $m$ are $u$ and $v$ in different connected components?

$(u,v) = e_i$

*if* $(u \text{ and } v \text{ are in different sets})$ {

$T \leftarrow T \cup \{e_i\}$

merge the sets containing $u$ and $v$

*

return $T$

}
Tie-breaking

To remove the assumption that all edge costs are distinct: perturb all edge costs by tiny amounts to break any ties.

Impact. Kruskal and Prim only interact with costs via pairwise comparisons. If perturbations are sufficiently small, MST with perturbed costs is MST with original costs. e.g., if all edge costs are integers, perturbing cost of edge $e_i$ by $i / n^2$

Implementation. Can handle arbitrarily small perturbations implicitly by instead breaking ties lexicographically, according to index.

```java
boolean less(i, j) {
    if (cost(e_i) < cost(e_j)) return true
    else if (cost(e_i) > cost(e_j)) return false
    else if (i < j) return true
    else return false
}
```


MST Algorithms: Theory

Deterministic comparison based algorithms.

- $O(m \log n)$  
  [Jarník, Prim, Dijkstra, Kruskal, Boruvka]
- $O(m \log \log n)$.  
  [Cheriton-Tarjan 1976, Yao 1975]
- $O(m \beta(m, n))$.  
  [Fredman-Tarjan 1987]
- $O(m \log \beta(m, n))$.  
  [Gabow-Galil-Spencer-Tarjan 1986]
- $O(m \alpha (m, n))$.  
  [Chazelle 2000]

Holy grail.  $O(m)$.

Notable.

- $O(m)$ randomized.  
  [Karger-Klein-Tarjan 1995]
- $O(m)$ verification.  
  [Dixon-Rauch-Tarjan 1992]

Euclidean.

- 2-d:  $O(n \log n)$.  
  compute MST of edges in Delaunay
- k-d:  $O(k n^2)$.  
  dense Prim
4.7 Clustering

Outbreak of cholera deaths in London in 1850s.
Reference: Nina Mishra, HP Labs
Clustering

Given a set $U$ of $n$ objects labeled $p_1, \ldots, p_n$, classify into coherent groups.

Distance function. Numeric value specifying "closeness" of two objects.

Fundamental problem. Divide into clusters so that points in different clusters are far apart.

- Routing in mobile ad hoc networks.
- Identify patterns in gene expression.
- Document categorization for web search.
- Similarity searching in medical image databases
- Skycat: cluster $10^9$ sky objects into stars, quasars, galaxies.
Dendrogram. Scientific visualization of hypothetical sequence of evolutionary events.

- Leaves = genes.
- Internal nodes = hypothetical ancestors.

Clustering of Maximum Spacing

**k-clustering.** Divide objects into k non-empty groups.

**Distance function.** Assume it satisfies several natural properties.
- $d(p_i, p_j) = 0$ iff $p_i = p_j$ (identity of indiscernible)
- $d(p_i, p_j) \geq 0$ (non-negativity)
- $d(p_i, p_j) = d(p_j, p_i)$ (symmetry)

**Spacing.** Min distance between any pair of points in different clusters.

**Clustering of maximum spacing.** Given an integer $k$, find a $k$-clustering of maximum spacing.

$k = 4$
Greedy Clustering Algorithm

Single-link k-clustering algorithm.
- Form a graph on the vertex set U, corresponding to n clusters.
- Find the closest pair of objects such that each object is in a different cluster, and add an edge between them.
- Repeat n-k times until there are exactly k clusters.

Key observation. This procedure is precisely Kruskal's algorithm (except that we stop with k (instead of 1) connected components).

Remark. Equivalent to finding an MST and deleting the k-1 most expensive edges.
**Theorem.** Let \( C^* \) denote the clustering \( C^*_1, \ldots, C^*_k \) formed by deleting the \( k-1 \) most expensive edges of an MST. \( C^* \) is a \( k \)-clustering of max spacing.

**Pf.** Let \( C \) denote some other clustering \( C_1, \ldots, C_k \).
- The **spacing** of \( C^* \) is the length \( d^* \) of the \((k-1)\)st most expensive edge.
- Let \( u, v \) be in the same cluster \( C^*_r \) in \( C^* \) but different clusters in \( C \).
- Some edge \((p, q)\) on \( u-v \) path in \( C^*_r \) spans two different clusters, say \( C_p \) and \( C_q \), in \( C \).
- All edges on \( u-v \) path have length \( \leq d^* \) since Kruskal chose them.
- Spacing of \( C \) is \( \leq d^* \) since \( p \) and \( q \) are in different clusters. ■
Chapter 4

Greedy Algorithms