

**McGill**APRIL 2013  
Final Examination

## FINAL EXAMINATION

**Computer Science COMP-547B**  
***Cryptography and Data Security***

29 APRIL 2013, 9h00

Examiner:	Prof. Claude Crépeau	Assoc Examiner:	Prof. Patrick Hayden
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**INSTRUCTIONS:**

- This examination is worth 50% of your final grade.
- The total of all questions is 105 points.
- Each question heading contains (in parenthesis) a list of values for each sub-questions.
- This is an **open book** exam. **All documentation is permitted.**
- Faculty standard calculator permitted only.
- The exam consists of 6 questions on 3 pages, title page included.

**Suggestion:****read all the questions and  
their values before you start.**

**Question 1. Perfect RSA? ( 10 + 10 = 20 points )**

Consider an RSA crypto-system with keys  $(N, e, d)$  as usual except that only  $N$  is publicly available.

- I) By definition, we know that  $m^{ed} \bmod N \equiv m$  for  $m$  s.t.  $\gcd(m, N) = 1$ . Using the Chinese remainder theorem show that  $m^{ed} \bmod N \equiv m$  for  $m$  s.t.  $\gcd(m, N) > 1$  as well.
- II) Assume Alice and Bob use  $(e, d)$  as the private encryption-decryption keys of an RSA crypto-system  $\bmod N$  for exactly one message  $m$ ,  $0 < m < N$ . Explain whether this one-time system is perfect according to Shannon's definition.

**Question 2. DDES ( 8 + 7 = 15 points )**

Consider the 128-bit block cipher DDES obtained by combining two instances of DES in a two-round Feistel network. The total key-size of this new cipher would be 112 bits.

- Let  $x$  be a 128-bit input and  $k$  be a 112-bit key. Give an explicit formula for the encryption and decryption functions of DDES.
- Discuss the pseudo-random nature of the permutation defined by DDES.

**Question 3. Rivest ( 10 + 10 = 20 points )**

Remember the construction by Rivest of a private-key crypto-system based on the existence of an arbitrary private-key authentication scheme.

- a) Show that the definition of security of the MAC is not sufficient for the resulting crypto-system to have undistinguishable encryptions in the presence of an eavesdropper.
- b) Define a stronger security notion for MACs such that the construction of Rivest yields a crypto-system with undistinguishable encryptions in the presence of an eavesdropper.

**Question 4. Number Theory vs Crypto ( 5 + 5 + 5 = 15 points )**

For each of the following Number Theoretical concepts, name a Cryptographic concept which is related and explain the relation.

- 1) Chinese remainder theorem.
- 2) Quadratic Residuosity.
- 3)  $\mathbb{F}_{2^k}$ , for  $k \geq 1$ .

**Question 5. Elgamal Details ( 10 + 10 = 20 points )**

Instantiate all the parameters of an Elgamal encryption scheme from a prime  $p=47$  and give me an encryption of  $m=10$ . Give me all the details of the crypto-system, taking into account all the implementation details seen in class.

(All your calculations can be done by hand.)

**Question 6. Merkle-Damgård... ( 8 + 7 = 15 points)**

In Construction 4.13 the size  $L$  of the input string  $x$  is such that  $L < 2^{l(n)}$ . It is very peculiar that if we hash a string  $x$  of length  $2^{l(n)} - 1$  using  $H^s$ , the time needed to hash the string is greater than the time needed to find a collision of  $h^s$  by a birthday attack. This seems to imply that hashing exponentially long strings is insecure.

- i) Explain why this is not contradicting the security statement (Theorem 4.14) that if  $h^s$  is collision-resistant then  $H^s$  is also collision-resistant.
- ii) Why do we still use an exponential bound ( $L < 2^{l(n)}$ ) in Construction 4.13 and not a polynomial bound such as  $L < l(n)^k$  ?